

# From Curves to Words and Back Again: Geometric Computation of Minimum-Area Homotopy

Hsien-Chih Chang, Brittany Terese Fasy,  
Bradley McCoy\*, David L. Millman, Carola Wenk

Montana State University  
August 1, 2023  
Montreal, Canada  
WADS

# Collaborator Appreciation



Hsien-Chih Chang



Brittany Fasy



David Millman

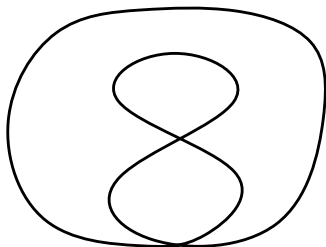


Carola Wenk

Supported by NSF grants CCF 2107434, DMS 1664858 and CCF 2046730

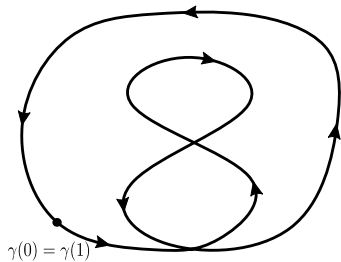
# Input to Our Problem

Generic Curves



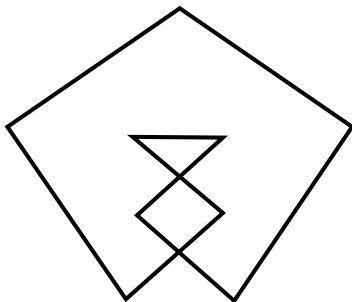
# Input to Our Problem

## Generic Curves



# Input to Our Problem

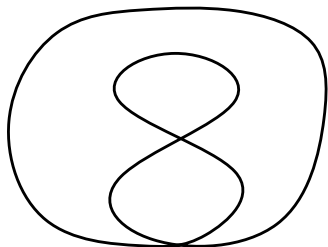
## Generic Curves



$n$  - number of segments,  $|F|$  - number of faces,  $|V|$  - number of self-intersections

# Input to Our Problem

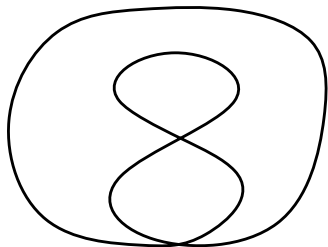
## Generic Curves



$n$  - number of segments,  $|F|$  - number of faces,  $|V|$  - number of self-intersections

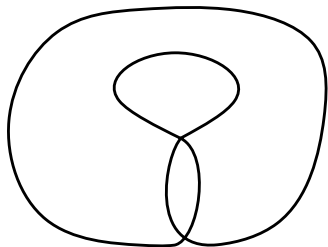
# Input to Our Problem

Minimum Area Homotopy



# Input to Our Problem

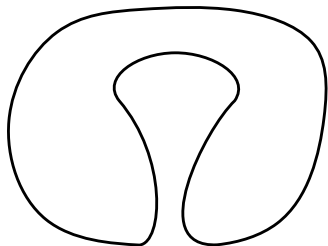
## Minimum Area Homotopy





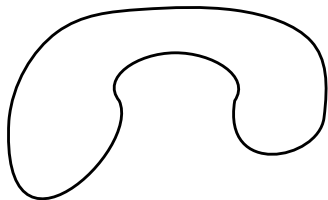
# Input to Our Problem

## Minimum Area Homotopy



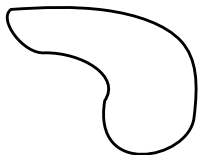
# Input to Our Problem

Minimum Area Homotopy



# Input to Our Problem

Minimum Area Homotopy



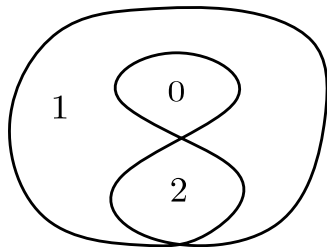
# Input to Our Problem

Minimum Area Homotopy



# Input to Our Problem

## Minimum Area Homotopy



# Agenda

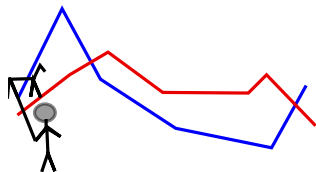
## Talk Outline

- Data Analysis
- Self-Overlapping Curves
- Minimum Area Homotopy
- Word Equivalence
- Mapping Class Groups
- An Open Inverse Problem

# Data Analysis Tool

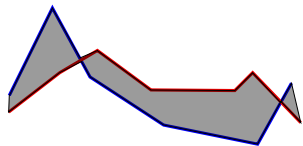
## Measure Curve Similarity

Fréchet Distance



$$\mathcal{O}(n^2 \log n)$$

Minimum Homotopy Area

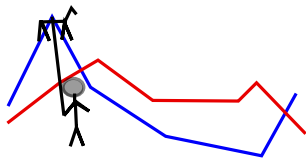


$$\mathcal{O}(n^2 + |F|^6)$$

# Data Analysis Tool

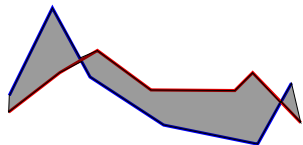
## Measure Curve Similarity

Fréchet Distance



$$\mathcal{O}(n^2 \log n)$$

Minimum Homotopy Area



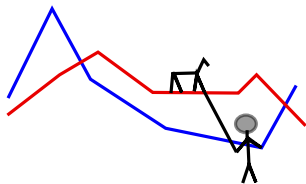
$$\mathcal{O}(n^2 + |F|^6)$$



# Data Analysis Tool

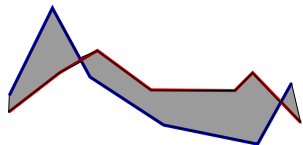
## Measure Curve Similarity

Fréchet Distance



$$\mathcal{O}(n^2 \log n)$$

Minimum Homotopy Area

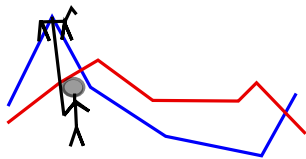


$$\mathcal{O}(n^2 + |F|^6)$$

# Data Analysis Tool

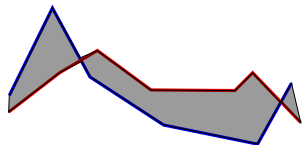
## Measure Curve Similarity

Fréchet Distance



$$\mathcal{O}(n^2 \log n)$$

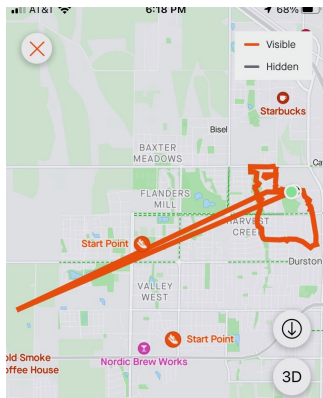
Minimum Homotopy Area



$$\mathcal{O}(n^2 + |F|^6)$$

# Data Analysis Tool

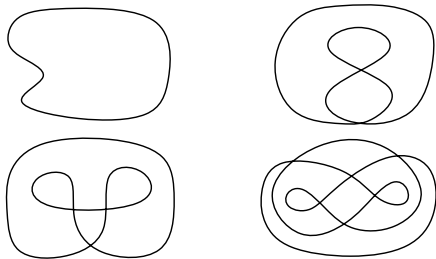
## Robust to Single Data Errors



# Is a Curve Self-Overlapping?

## Definition

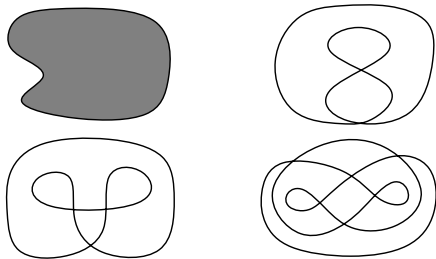
A curve  $\gamma$  is **self-overlapping** if there exists an immersion  $F : D \rightarrow \mathbb{R}^2$  with  $\gamma = F|_{\partial(D)}$ .



# Is a Curve Self-Overlapping?

## Definition

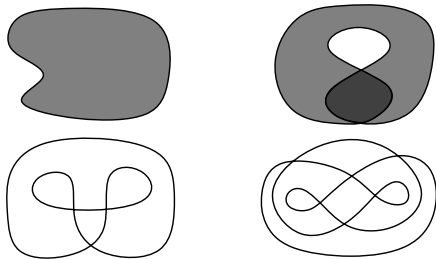
A curve  $\gamma$  is **self-overlapping** if there exists an immersion  $F : D \rightarrow \mathbb{R}^2$  with  $\gamma = F|_{\partial(D)}$ .



# Is a Curve Self-Overlapping?

## Definition

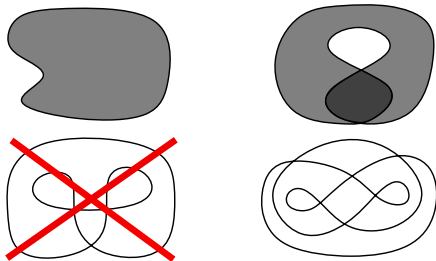
A curve  $\gamma$  is **self-overlapping** if there exists an immersion  $F : D \rightarrow \mathbb{R}^2$  with  $\gamma = F|_{\partial(D)}$ .



# Is a Curve Self-Overlapping?

## Definition

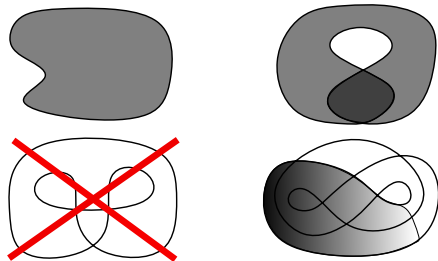
A curve  $\gamma$  is **self-overlapping** if there exists an immersion  $F : D \rightarrow \mathbb{R}^2$  with  $\gamma = F|_{\partial(D)}$ .



# Is a Curve Self-Overlapping?

## Definition

A curve  $\gamma$  is **self-overlapping** if there exists an immersion  $F : D \rightarrow \mathbb{R}^2$  with  $\gamma = F|_{\partial(D)}$ .

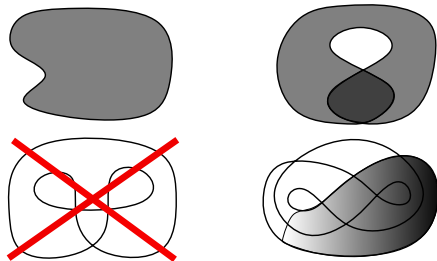




# Is a Curve Self-Overlapping?

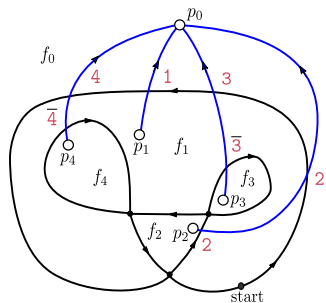
## Definition

A curve  $\gamma$  is **self-overlapping** if there exists an immersion  $F : D \rightarrow \mathbb{R}^2$  with  $\gamma = F|_{\partial(D)}$ .



# Blank's Algorithm

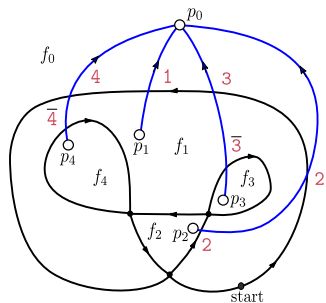
## Blank Word Construction



$$W = []$$

# Blank's Algorithm

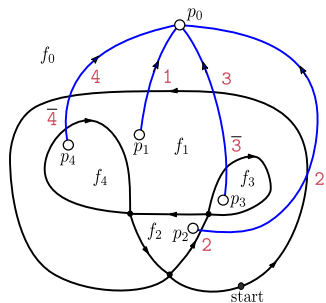
## Blank Word Construction



$$W = [2]$$

# Blank's Algorithm

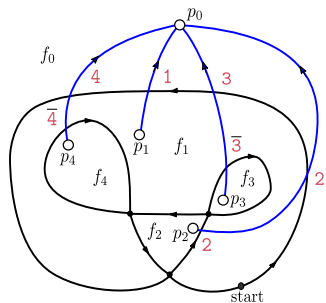
## Blank Word Construction



$$W = [23]$$

# Blank's Algorithm

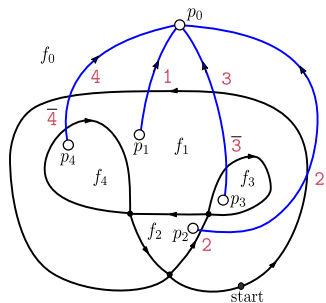
## Blank Word Construction



$$W = [231]$$

# Blank's Algorithm

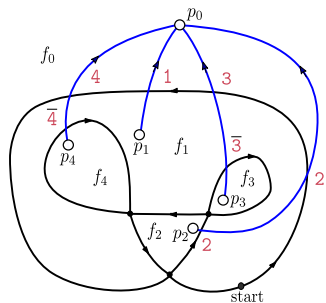
## Blank Word Construction



$$W = [2314]$$

# Blank's Algorithm

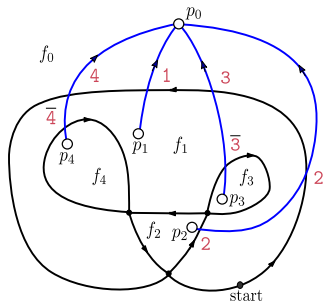
## Blank Word Construction



$$W = [23142]$$

# Blank's Algorithm

## Blank Word Construction

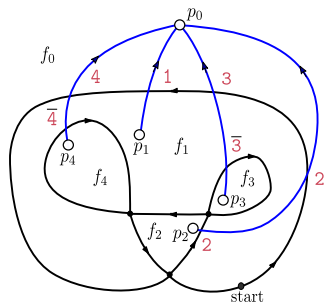


$$W = [23142\bar{3}]$$



# Blank's Algorithm

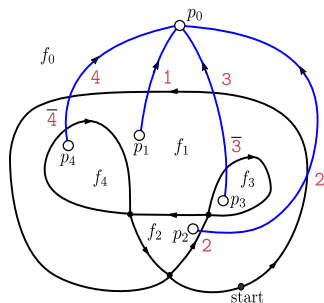
## Blank Word Construction



$$W = [23142\bar{3}\bar{4}]$$

# Blank's Algorithm

## Blank Cuts and Groupings

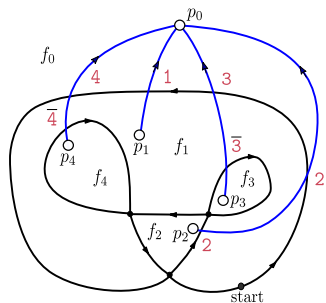


$$w = [23142\bar{3}\bar{4}]$$

- Positive subword  $\sigma = f_1 f_2 \dots f_k$  where each  $f_i$  is positive

# Blank's Algorithm

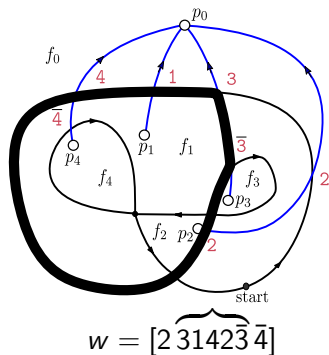
## Blank Cuts and Groupings



- Positive subword  $\sigma = f_1 f_2 \dots f_k$  where each  $f_i$  is positive
- Positive pairing  $(f, \bar{f})$ ,  $w = [fp\bar{f}w']$   $p$  positive word

# Blank's Algorithm

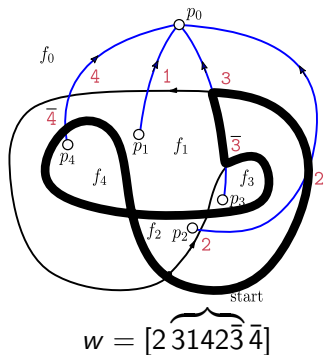
## Blank Cuts and Groupings



- Positive subword  $\sigma = f_1 f_2 \dots f_k$  where each  $f_i$  is positive
- Positive pairing  $(f, \bar{f})$ ,  $w = [fp\bar{f}w']$   $p$  positive word
- Blank cut - replace positive pairing with identity

# Blank's Algorithm

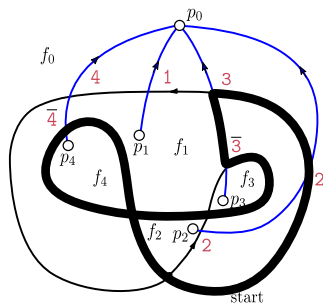
## Blank Cuts and Groupings



- Positive subword  $\sigma = f_1 f_2 \dots f_k$  where each  $f_i$  is positive
- Positive pairing  $(f, \bar{f})$ ,  $w = [fp\bar{f}w']$   $p$  positive word
- Blank cut - replace positive pairing with identity

# Blank's Algorithm

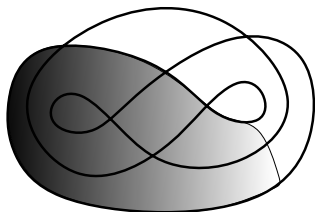
## Blank Cuts and Groupings



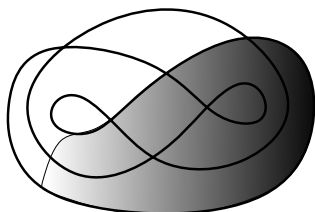
- *Positive* subword  $\sigma = f_1 f_2 \dots f_k$  where each  $f_i$  is positive
- *Positive pairing*  $(f, \bar{f})$ ,  $w = [fp\bar{f}w']$   $p$  positive word
- *Blank cut* - replace positive pairing with identity
- *Groupable* - remove positive pairings until positive word

# Blank's Algorithm

Count the Number of Extensions



$$w = a_2 a_5 a_1 a_3 a_4 a_6 a_5^{-1} a_3 a_4 a_6 a_7 a_2 a_5 a_3 a_6^{-1}$$

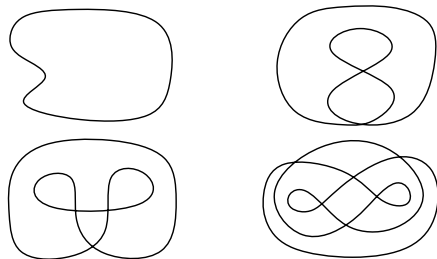


$$w = a_2 a_5 a_1 a_3 a_4 a_6 a_5^{-1} a_3 a_4 a_6 a_7 a_2 a_5 a_3 a_6^{-1}$$

# Examples

## Definition

A **minimum area homotopy** is a continuous deformation from a curve  $\gamma$  to a point that sweeps over the minimum area in the plane.

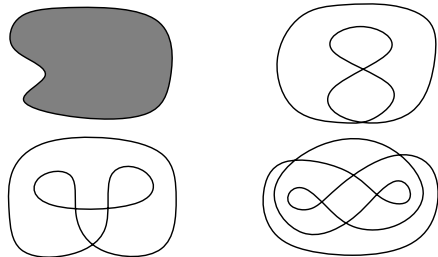




# Examples

## Definition

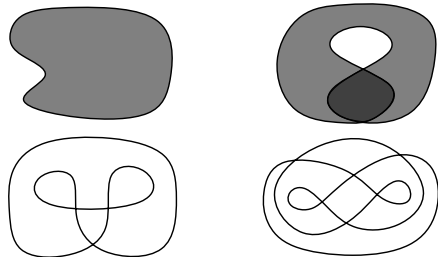
A **minimum area homotopy** is a continuous deformation from a curve  $\gamma$  to a point that sweeps over the minimum area in the plane.



# Examples

## Definition

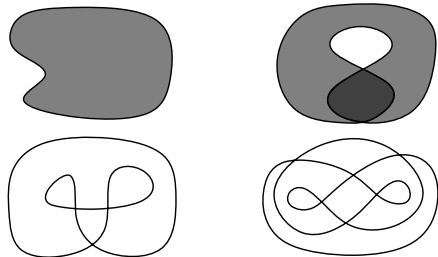
A **minimum area homotopy** is a continuous deformation from a curve  $\gamma$  to a point that sweeps over the minimum area in the plane.



# Examples

## Definition

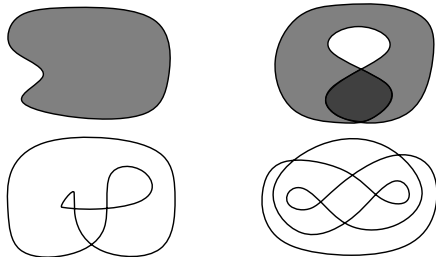
A **minimum area homotopy** is a continuous deformation from a curve  $\gamma$  to a point that sweeps over the minimum area in the plane.



# Examples

## Definition

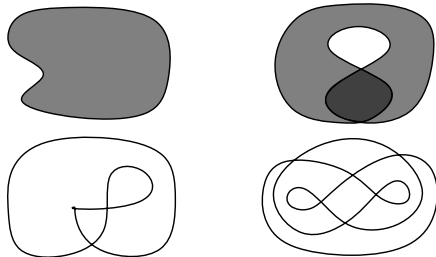
A **minimum area homotopy** is a continuous deformation from a curve  $\gamma$  to a point that sweeps over the minimum area in the plane.



# Examples

## Definition

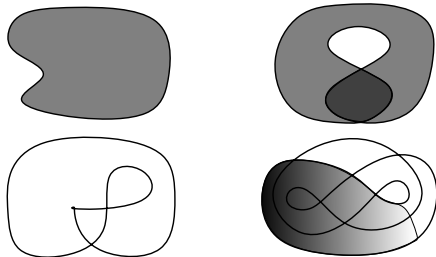
A **minimum area homotopy** is a continuous deformation from a curve  $\gamma$  to a point that sweeps over the minimum area in the plane.



# Examples

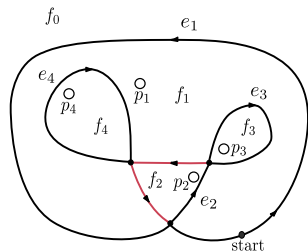
## Definition

A **minimum area homotopy** is a continuous deformation from a curve  $\gamma$  to a point that sweeps over the minimum area in the plane.



# Nie's Algorithm

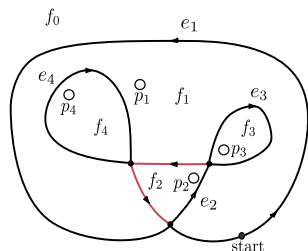
## Words From Boundaries



- Choose a spanning tree

# Nie's Algorithm

## Words From Boundaries

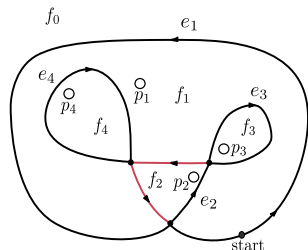


- Choose a spanning tree
- Write each face as a boundary of edges
  - $\partial(f_1) = e_3 \bar{e}_2 e_1 e_4$
  - $\partial(f_2) = e_2$
  - $\partial(f_3) = \bar{e}_3$
  - $\partial(f_4) = \bar{e}_4$



# Nie's Algorithm

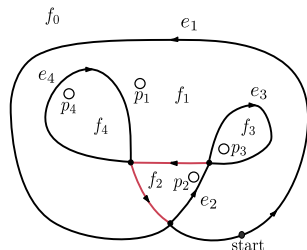
## Words From Boundaries



- Choose a spanning tree
- Write each face as a boundary of edges
  - $\partial(f_1) = e_3 \bar{e}_2 e_1 e_4$
  - $\partial(f_2) = e_2$
  - $\partial(f_3) = \bar{e}_3$
  - $\partial(f_4) = \bar{e}_4$
- Write the curve as  $\gamma = e_1 e_2 e_3 e_4$

# Nie's Algorithm

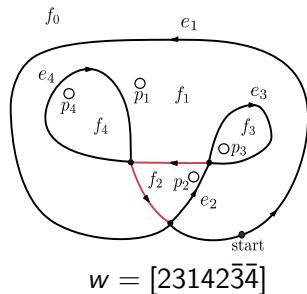
## Words From Boundaries



- Choose a spanning tree
- Write each face as a boundary of edges
  - $\partial(f_1) = e_3 \bar{e}_2 e_1 e_4$
  - $\partial(f_2) = e_2$
  - $\partial(f_3) = \bar{e}_3$
  - $\partial(f_4) = \bar{e}_4$
- Write the curve as  $\gamma = e_1 e_2 e_3 e_4$
- $\gamma = \partial(f_2) \partial(f_3) \partial(f_1) \partial(f_4) \partial(f_2) \partial(\bar{f}_3) \partial(\bar{f}_4)$

# Nie's Algorithm

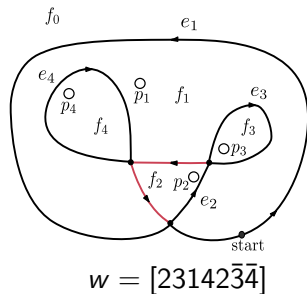
## Words From Boundaries



- Folding - set of unlinked pairings

# Nie's Algorithm

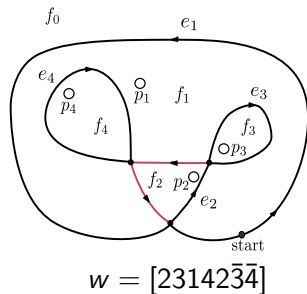
## Words From Boundaries



- Folding - set of unlinked pairings
- Weighted cancellation norm

# Nie's Algorithm

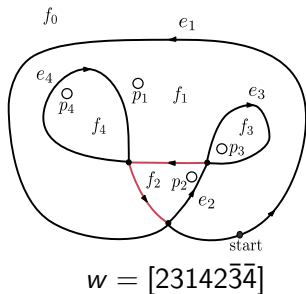
## Words From Boundaries



- Folding - set of unlinked pairings
- Weighted cancellation norm
- $\|w\| := \min_{\mathcal{F}} \sum_i \text{Area}(f_i)$

# Nie's Algorithm

## Words From Boundaries

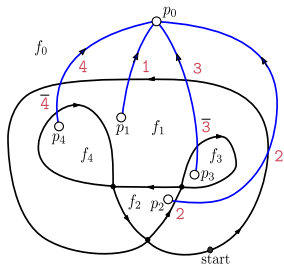


- Folding - set of unlinked pairings
- Weighted cancellation norm
- $\|w\| := \min_{\mathcal{F}} \sum_i \text{Area}(f_i)$
- Dynamic program - think matrix chain multiplication

# Word Equivalence

Where Did We Make Choices?

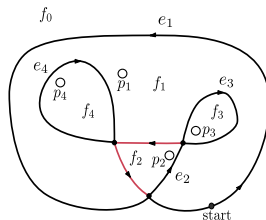
Blank's Word



$$w = [12342\bar{3}\bar{4}]$$

Nie's Word

$$\partial(f_1) = e_3 \bar{e}_2 e_1 e_4$$

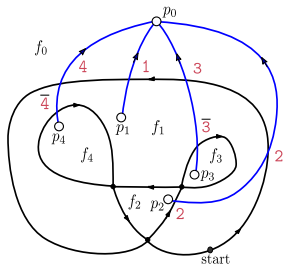


$$w = [12342\bar{3}\bar{4}]$$

# Word Equivalence

Where Did We Make Choices?

Blank's Word

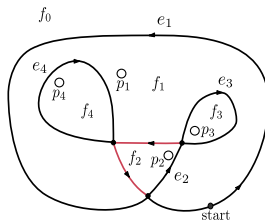


$$w = [12342\bar{3}\bar{4}]$$

How to draw the cables.

Nie's Word

$$\partial(f_1) = e_3 \bar{e}_2 e_1 e_4$$



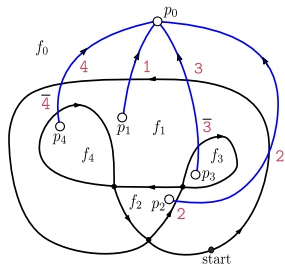
$$w = [12342\bar{3}\bar{4}]$$



# Word Equivalence

Where Did We Make Choices?

Blank's Word

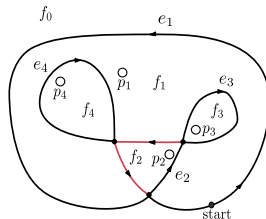


$$w = [12342\bar{3}\bar{4}]$$

How to draw the cables.

Nie's Word

$$\partial(f_1) = e_3 \bar{e}_2 e_1 e_4$$



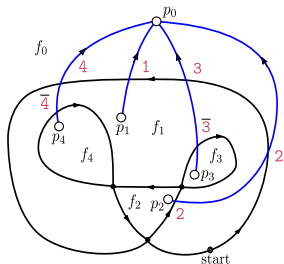
$$w = [12342\bar{3}\bar{4}]$$

Spanning tree.

# Word Equivalence

Where Did We Make Choices?

Blank's Word



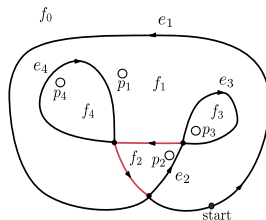
$$w = [12342\bar{3}\bar{4}]$$

How to draw the cables.

Order around  $p_0$ .

Nie's Word

$$\partial(f_1) = e_3 \bar{e}_2 e_1 e_4$$



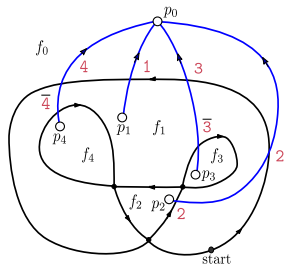
$$w = [12342\bar{3}\bar{4}]$$

Spanning tree.

# Word Equivalence

Where Did We Make Choices?

Blank's Word



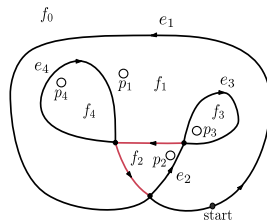
$$w = [12342\bar{3}\bar{4}]$$

How to draw the cables.

Order around  $p_0$ .

Nie's Word

$$\partial(f_1) = e_3 \bar{e}_2 e_1 e_4$$



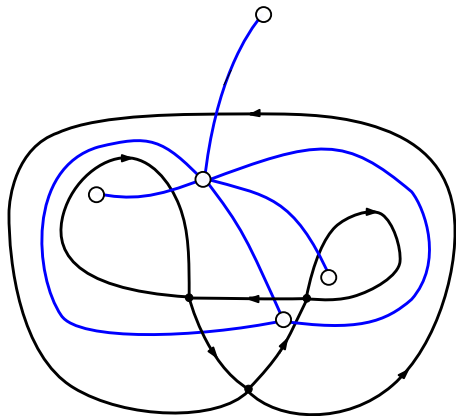
$$w = [12342\bar{3}\bar{4}]$$

Spanning tree.

Cyclic order of boundary.

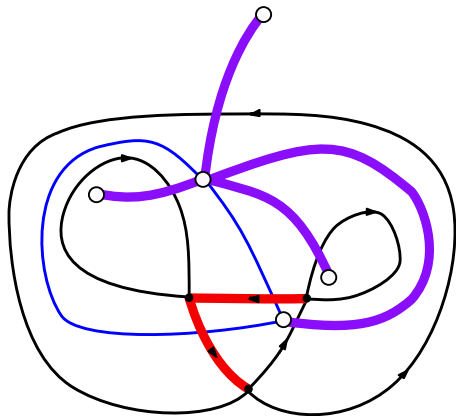
# Word Equivalence

## Tree Co-Tree Decomposition



# Word Equivalence

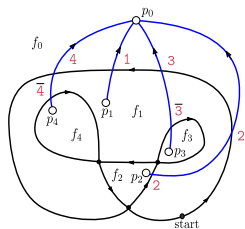
## Tree Co-Tree Decomposition



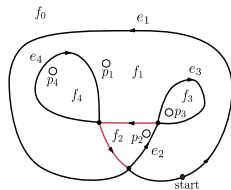
# Word Equivalence

Where Did We Make Choices?

Blank's Word



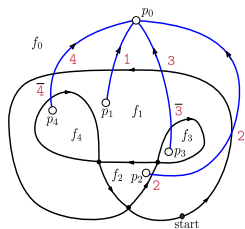
Nie's Word



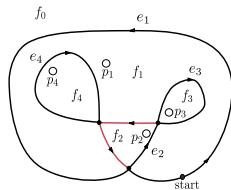
# Word Equivalence

Where Did We Make Choices?

Blank's Word



Nie's Word

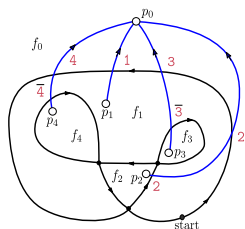


Where to place cable one.

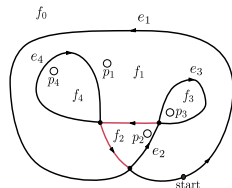
# Word Equivalence

Where Did We Make Choices?

Blank's Word



Nie's Word



$$\partial(f_1) = e_4 e_3 \bar{e}_2 e_1 \sim e_1 e_4 e_3 \bar{e}_2$$

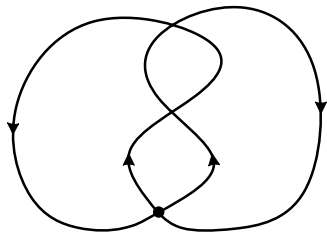
Where to place cable one.



# Foldings $\iff$ Self-Overlapping Decompositions

Theorem (Fasy, Karakoç, Wenk - 2017)

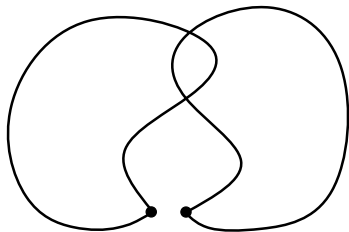
*Any curve has a self-overlapping decomposition whose area is minimum over all null-homotopies*



# Foldings $\iff$ Self-Overlapping Decompositions

Theorem (Fasy, Karakoç, Wenk - 2017)

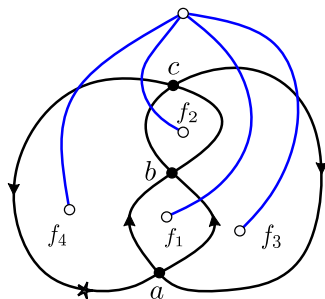
*Any curve has a self-overlapping decomposition whose area is minimum over all null-homotopies*



# Foldings $\iff$ Self-Overlapping Decompositions

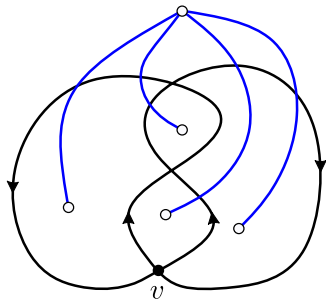
Theorem (Fasy, Karakoç, Wenk - 2017)

*Any curve has a self-overlapping decomposition whose area is minimum over all null-homotopies*

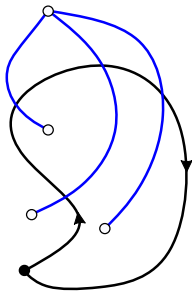


$$w = [a1b\bar{2}c\bar{1}\bar{3}abc24].$$

# Possible Issue



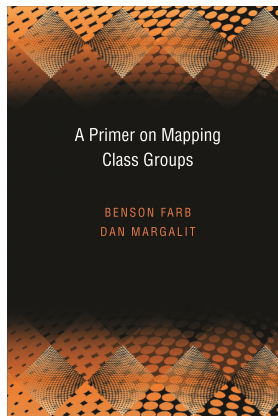
# Possible Issue



# Cable Independence

## Mapping Class Groups

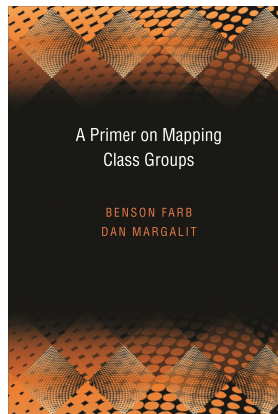
- surface with genus  $g$  and  $n$  punctures  $S_{g,n}$



# Cable Independence

## Mapping Class Groups

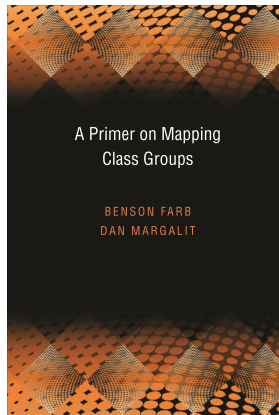
- surface with genus  $g$  and  $n$  punctures  $S_{g,n}$
- group of orientation-preserving diffeomorphisms  $\text{Diffeo}^+(S_{g,n})$



# Cable Independence

## Mapping Class Groups

- surface with genus  $g$  and  $n$  punctures  $S_{g,n}$
- group of orientation-preserving diffeomorphisms  $\text{Diffeo}^+(S_{g,n})$
- equivalence relation  $\sim$  on  $\text{Diffeo}$   
 $\phi \sim \psi$  if  $\phi$  and  $\psi$  isotopic

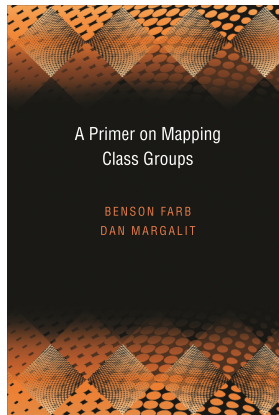




# Cable Independence

## Mapping Class Groups

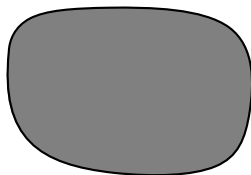
- surface with genus  $g$  and  $n$  punctures  $S_{g,n}$
- group of orientation-preserving diffeomorphisms  $\text{Diffeo}^+(S_{g,n})$
- equivalence relation  $\sim$  on  $\text{Diffeo}^+$   
 $\phi \sim \psi$  if  $\phi$  and  $\psi$  isotopic
- the mapping class group  
 $\text{MCG}(S_{g,n}) = \text{Diffeo}^+(S_{g,n}) / \sim$ .



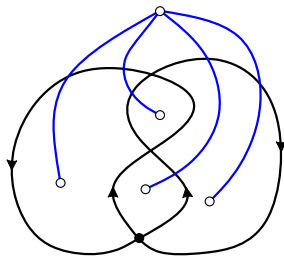
# Cable Independence

## Mapping Class Groups

- surface with genus  $g$  and  $n$  punctures  $S_{g,n}$
- group of orientation-preserving diffeomorphisms  $\text{Diffeo}^+(S_{g,n})$
- equivalence relation  $\sim$  on  $\text{Diffeo}^+$   
 $\phi \sim \psi$  if  $\phi$  and  $\psi$  isotopic
- the mapping class group  
 $\text{MCG}(S_{g,n}) = \text{Diffeo}^+(S_{g,n}) / \sim$ .

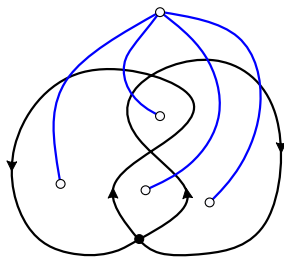


# Cable Independence



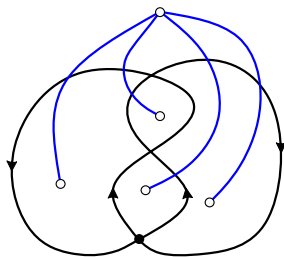
- our surface - punctured disk

# Cable Independence



- our surface - punctured disk
- cable - nonseparating puncture-to-puncture arc

# Cable Independence



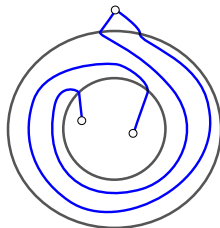
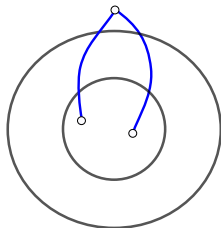
- our surface - punctured disk
- cable - nonseparating puncture-to-puncture arc
- pure mapping class group - punctures fixed

# Cable Independence

## Dehn Twists

### Theorem (9.3 Farb, Margalit)

*The pure mapping class group of  $D_n$  is generated by Dehn twists about the set of simple closed curves that surround exactly two punctures.*

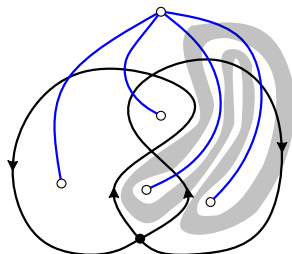
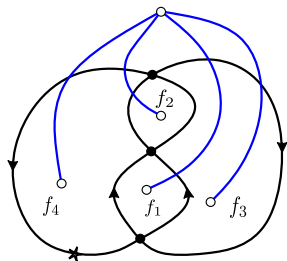


# Cable Independence

## Dehn Twists

### Theorem (9.3 Farb, Margalit)

*The pure mapping class group of  $D_n$  is generated by Dehn twists about the set of simple closed curves that surround exactly two punctures.*



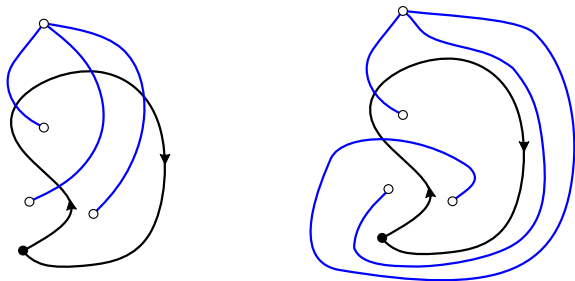
$$[1\overline{21}3\overline{24}] \rightarrow [3\overline{11}1\overline{32}3\overline{11}1\overline{33}1\overline{31}3\overline{24}]$$

# Cable Independence

## Dehn Twists

### Theorem (9.3 Farb, Margalit)

*The pure mapping class group of  $D_n$  is generated by Dehn twists about the set of simple closed curves that surround exactly two punctures.*





# The Results

## Lemma

*Let  $\gamma$  be a curve. For each folding  $F$  there exists a null-homotopy of  $\gamma$  with area equal to the area of  $F$ .*

# The Results

## Lemma

*Let  $\gamma$  be a curve. For each folding  $F$  there exists a null-homotopy of  $\gamma$  with area equal to the area of  $F$ .*

## Theorem

*Given a self-overlapping decomposition  $\Gamma$  of  $\gamma$ , there exists a folding  $F$  of  $w$  whose area is  $\text{Area}_\Gamma(\gamma)$ .*

# The Results

## Lemma

*Let  $\gamma$  be a curve. For each folding  $F$  there exists a null-homotopy of  $\gamma$  with area equal to the area of  $F$ .*

## Theorem

*Given a self-overlapping decomposition  $\Gamma$  of  $\gamma$ , there exists a folding  $F$  of  $w$  whose area is  $\text{Area}_\Gamma(\gamma)$ .*

## Corollary

*New proof of correctness.*

# The Results

## Theorem

*Given a maximal folding  $F$  of  $w$ , there is a self-overlapping decomposition of  $\gamma$  whose area is equal to the area induced by the folding  $F$ .*

# The Results

## Theorem

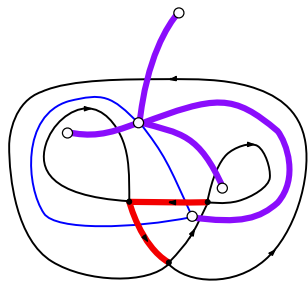
*Given a maximal folding  $F$  of  $w$ , there is a self-overlapping decomposition of  $\gamma$  whose area is equal to the area induced by the folding  $F$ .*

## Corollary

*Polynomial time minimum area self-overlapping decomposition.*

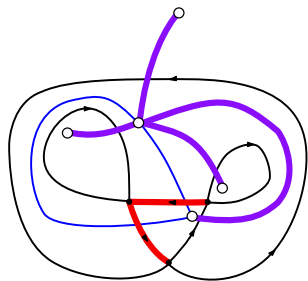
# Highlight Cool Ideas

## Cut Cycle Duality

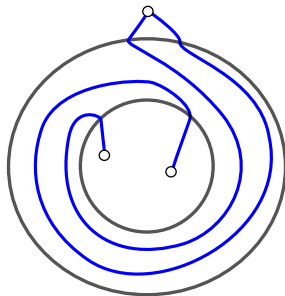


# Highlight Cool Ideas

## Cut Cycle Duality



## Dehn Twists



# An Open Problem

## Which Words are Curves?

Given a word, construct a curve with word equal to the given word or say no curve exists.

$$w_1 = [12]$$

$$w_2 = [122]$$

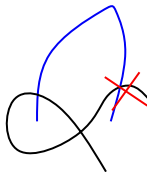


# An Open Problem

## Which Words are Curves?

Given a word, construct a curve with word equal to the given word or say no curve exists.

$$w_1 = [12]$$



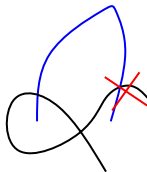
$$w_2 = [122]$$

# An Open Problem

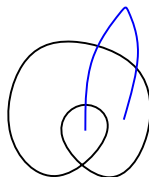
## Which Words are Curves?

Given a word, construct a curve with word equal to the given word or say no curve exists.

$$w_1 = [12]$$

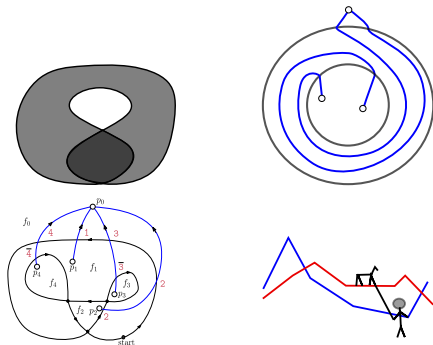


$$w_2 = [122]$$



# Questions?

Thank You



[bradleymccoy@montana.edu](mailto:bradleymccoy@montana.edu)